



Information Coding / Computer Graphics, ISY, LiTH

Collision detection and handling

A major problem in many real-time animations.

Even modern high-budget games have problems!



2D collision detection

Even in 2D, collisions is a non-trivial problem.

- **Rectanges - easy**
- **Enclosing convex polygon(s)**
 - **Pixel-level testing**



2D collision detection with rectangles

Testing with the bounding box: often unsatisfactory. Too crude approximation!



Smaller rectangles make decent compromises.





Special cases

When using pseudo-3D, sprites in perspective, the shape outline is not usable.

Simple approach: Modified box:



Wrong

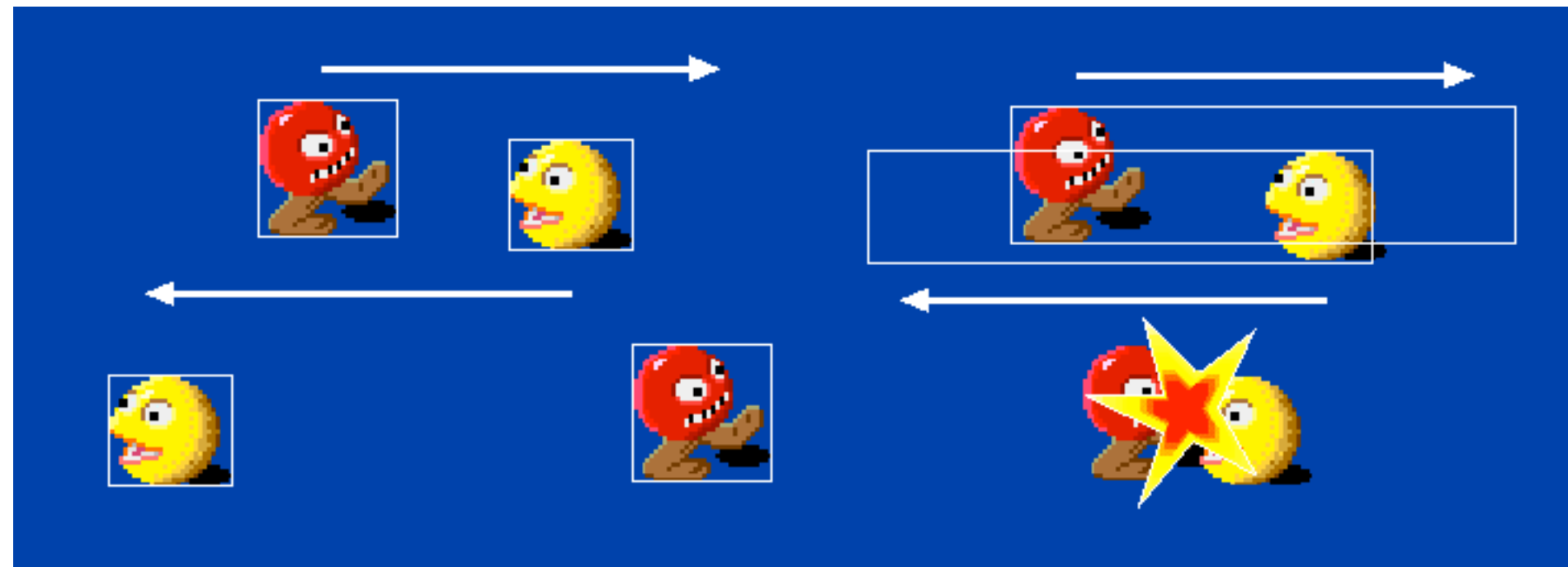


Right



Moving objects

Fast moving objects may pass right through each other.



Wrong

Right

An example of *temporal aliasing*



Fast-moving objects

- **Test with elongated shapes (sweeping)**
- **Test several times along the trajectory (multisampling)**



Polyhedra-polyhedra collisions

**Naive method: Check all polygons in all
objects against all polygons in all
objects!**



Polyhedra-polyhedra collisions

Practical methods:

- 1) Separate into broad phase and narrow phase**
- 2) Object hierarchies**



Broad phase - narrow phase

Really three phases:

Global phase
Bounding shapes phase
Narrow phase

) **Broad phase**



Bounding shapes phase

Make simple checks to see if objects are so close that we should test in detail!

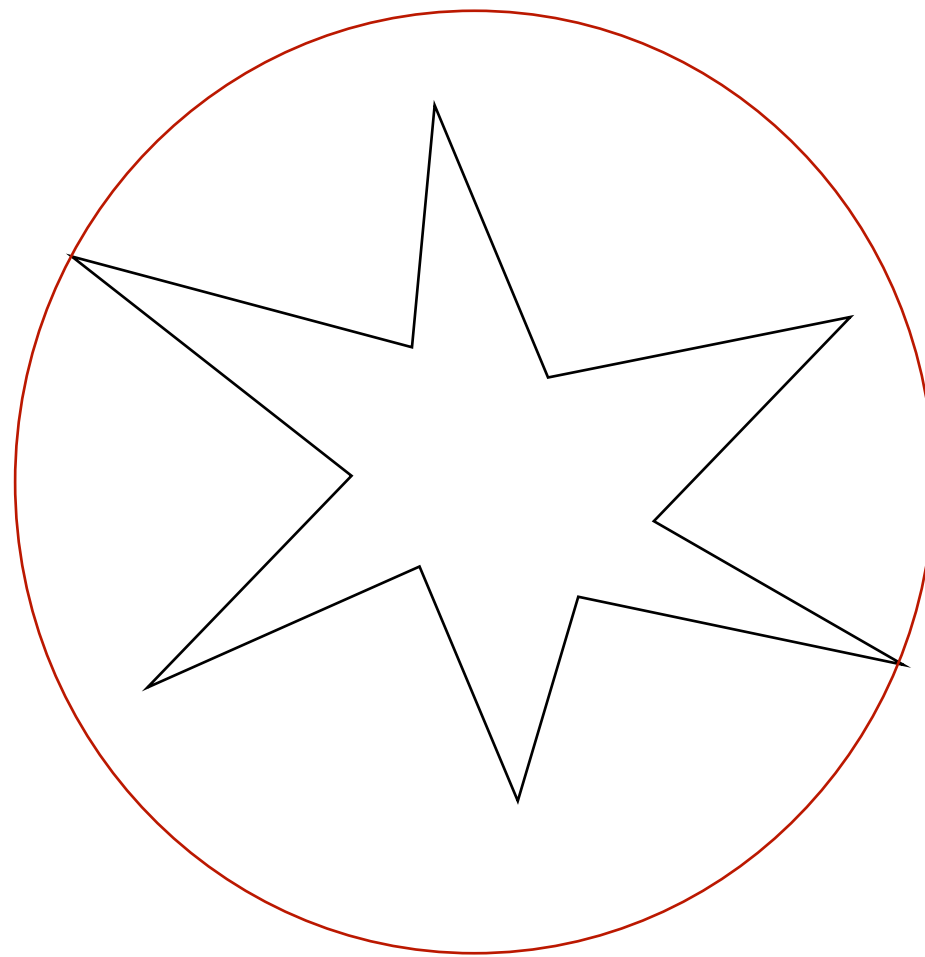
Use simple bounding shapes!

Typical shapes:

- **Sphere**
- **AABB (axis aligned bounding box)**
 - **OBB (oriented bounding box)**

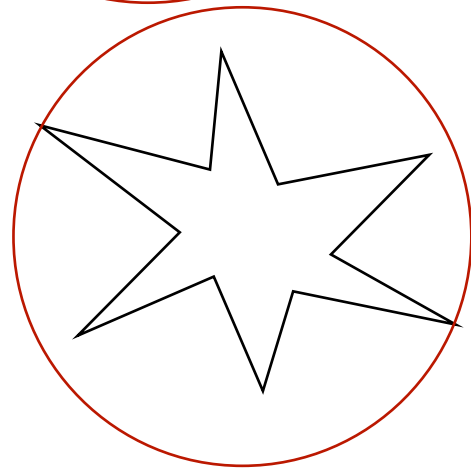
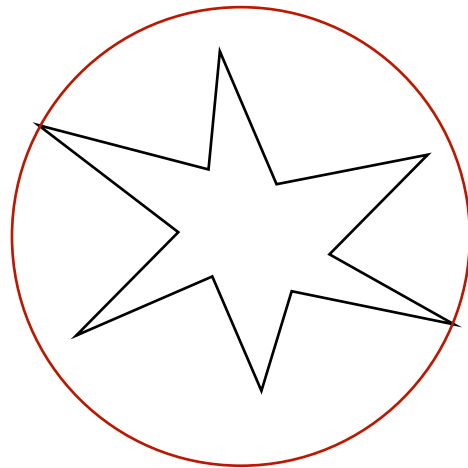


Sphere





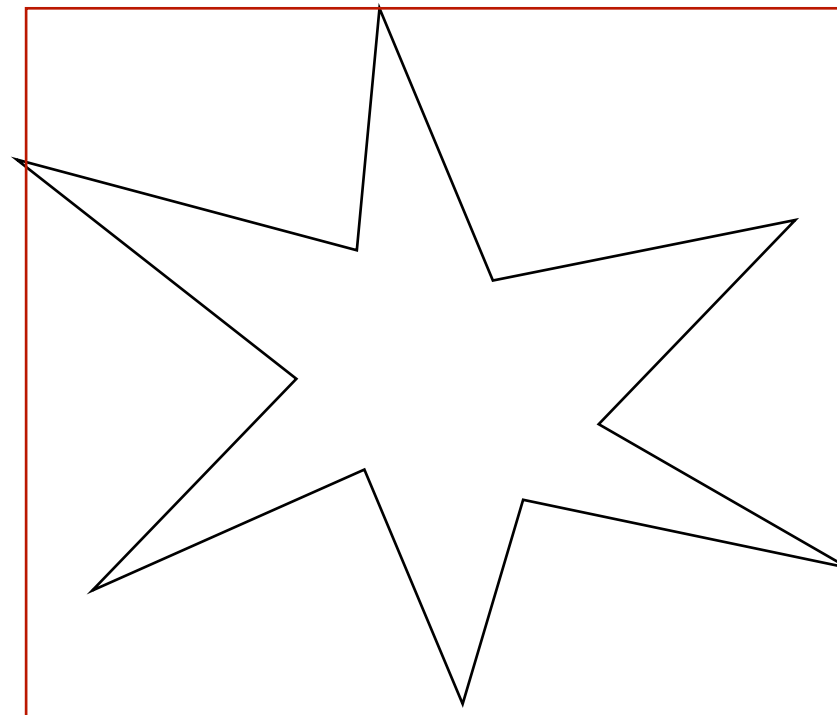
Sphere



- **Very simple tests. $d^2 < (r_1 + r_2)^2$**
- **Calculate radius once and for all. Never changes for rigid objects. Loop through all vertices, pick biggest distance to center.**
- **Rotation allowed!**
- **Good fit for compact objects.**
- **No flat surfaces, object can not rest on each other.**



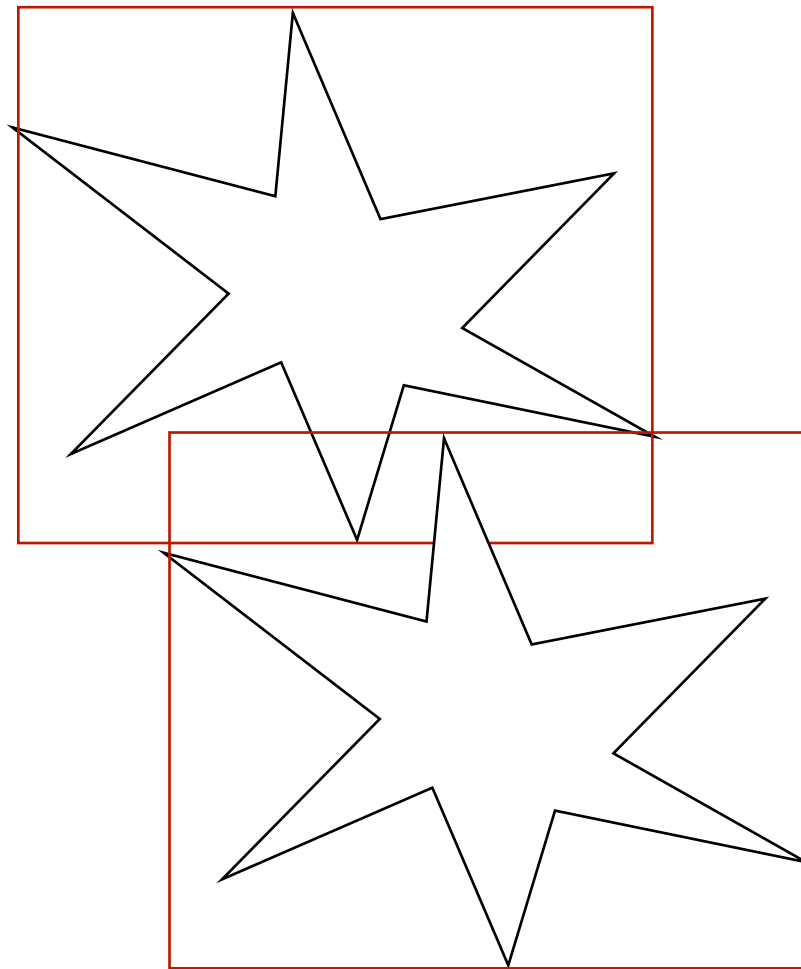
AABB



Axis aligned bounding box



AABB

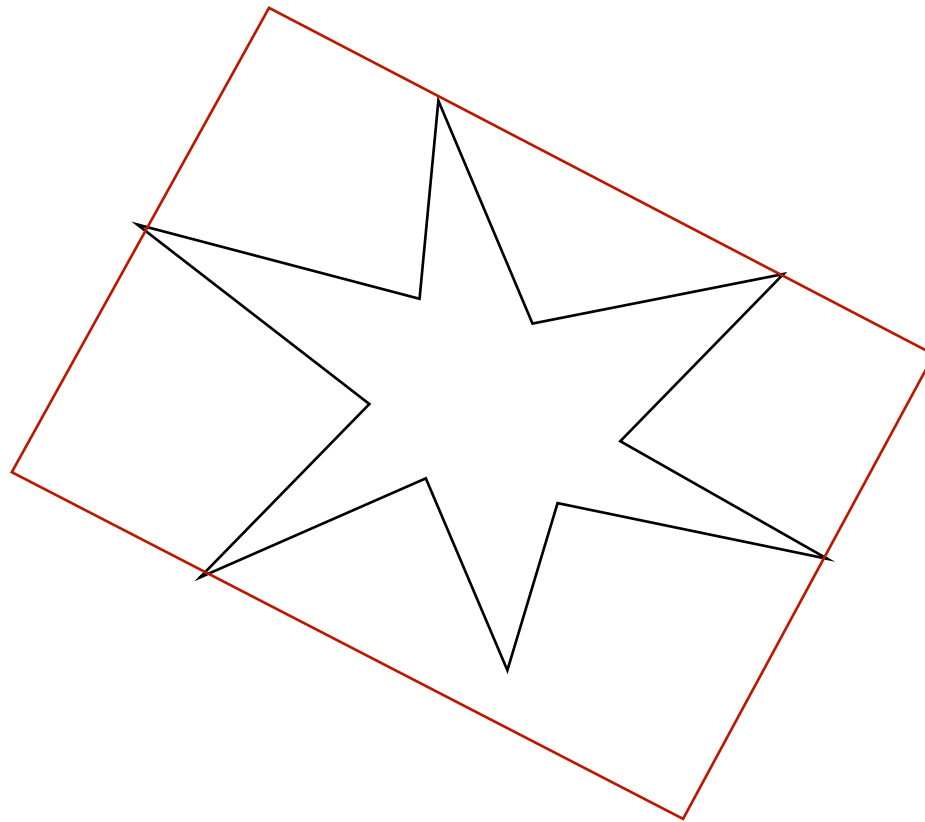


- **Very simple tests. 6 tests, like the 2D case.**
- **Calculate by finding max and min along each axis.**
- **Can't be rotated freely!**
 - 1) **Recalculate on rotation.**
 - 2) **Make new AABB from rotated vertices of AABB.**
- **Not very good fit on many objects.**



OBB

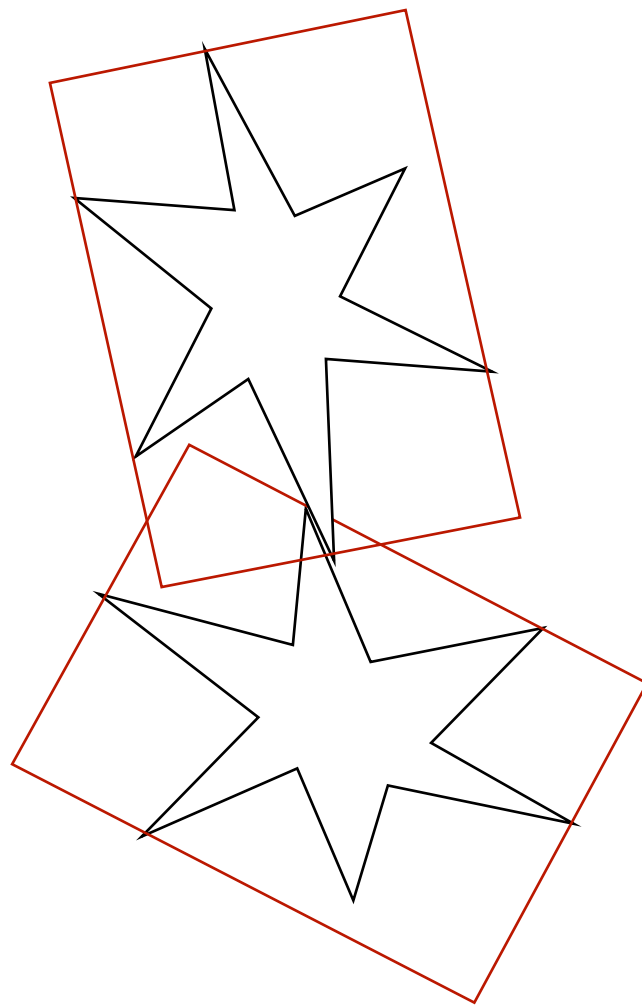
Oriented bounding box





OBB

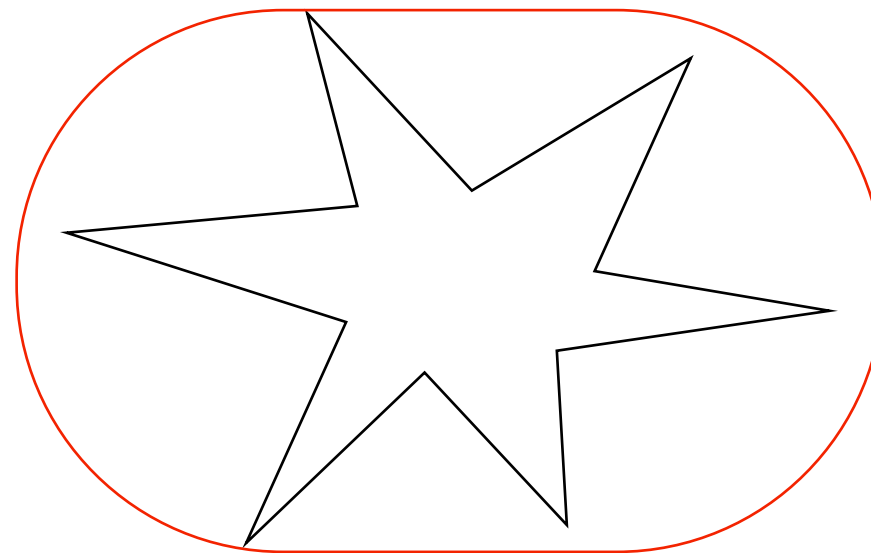
Oriented bounding box



- **Good fit, you can find smallest possible box. Few false hits.**
- **Complicated tests.**
- **Rotation allowed.**



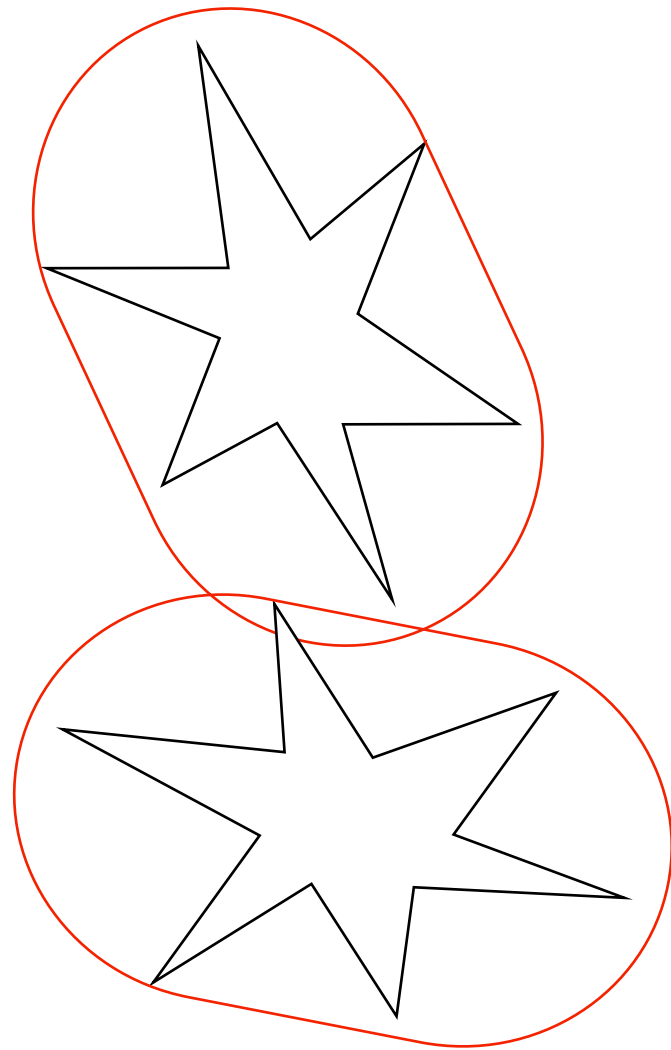
Capsule





Capsule

Cylinder + 2 half spheres

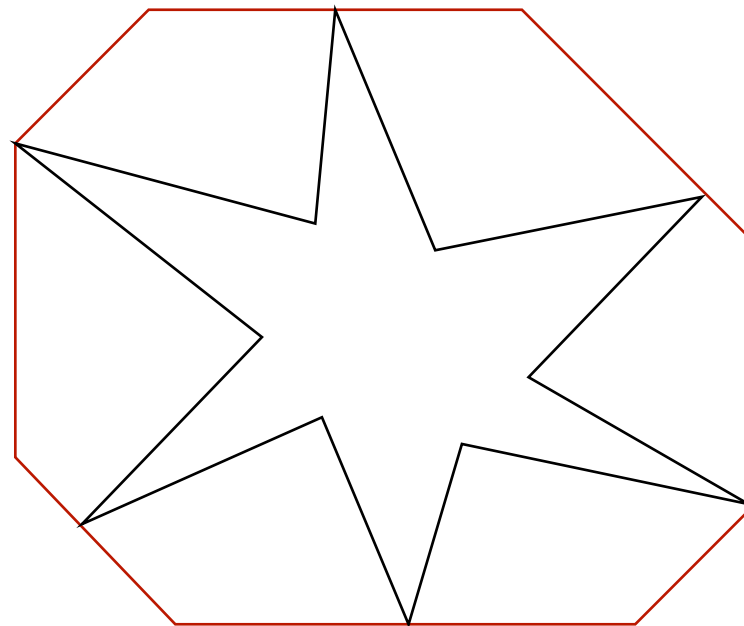


- **Good fit, you can find smallest possible box. Few false hits.**
- **Fairly complicated tests - but easier than OBB.**
- **Rotation allowed.**



k-DOP

Discrete oriented polytope with k sides



Like AABB but you can cut off corners and sides
Fairly simple tests
Fewer false hits

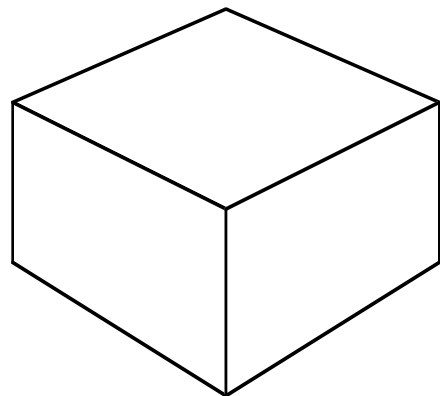


k-DOP

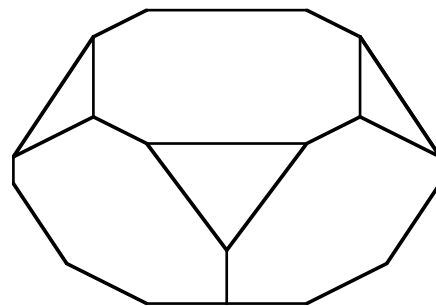
Discrete oriented polytope with k sides

Four kinds:

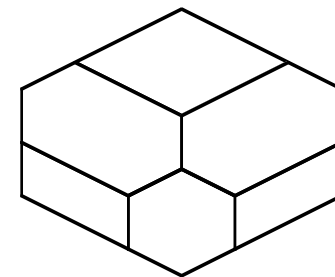
- 6-DOP (= AABB)**
- 14-DOP (corners)**
- 18-DOP (sides)**
- 26-DOP (sides and corners)**



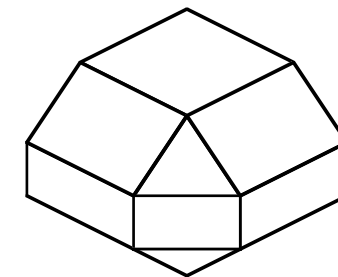
6-DOP



14-DOP



18-DOP



26-DOP



Balancing the broad and narrow phases

$$T = N_v C_v + N_p C_p$$

Simple broad phase - many false hits - high N_p .

**Elaborate broad phase - few false hits - lower N_p
but high C_v .**



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Narrow phase

Separating Axis Theorem (SAT)

Containment/intersection test

Volume intersection

Advanced methods: GJK



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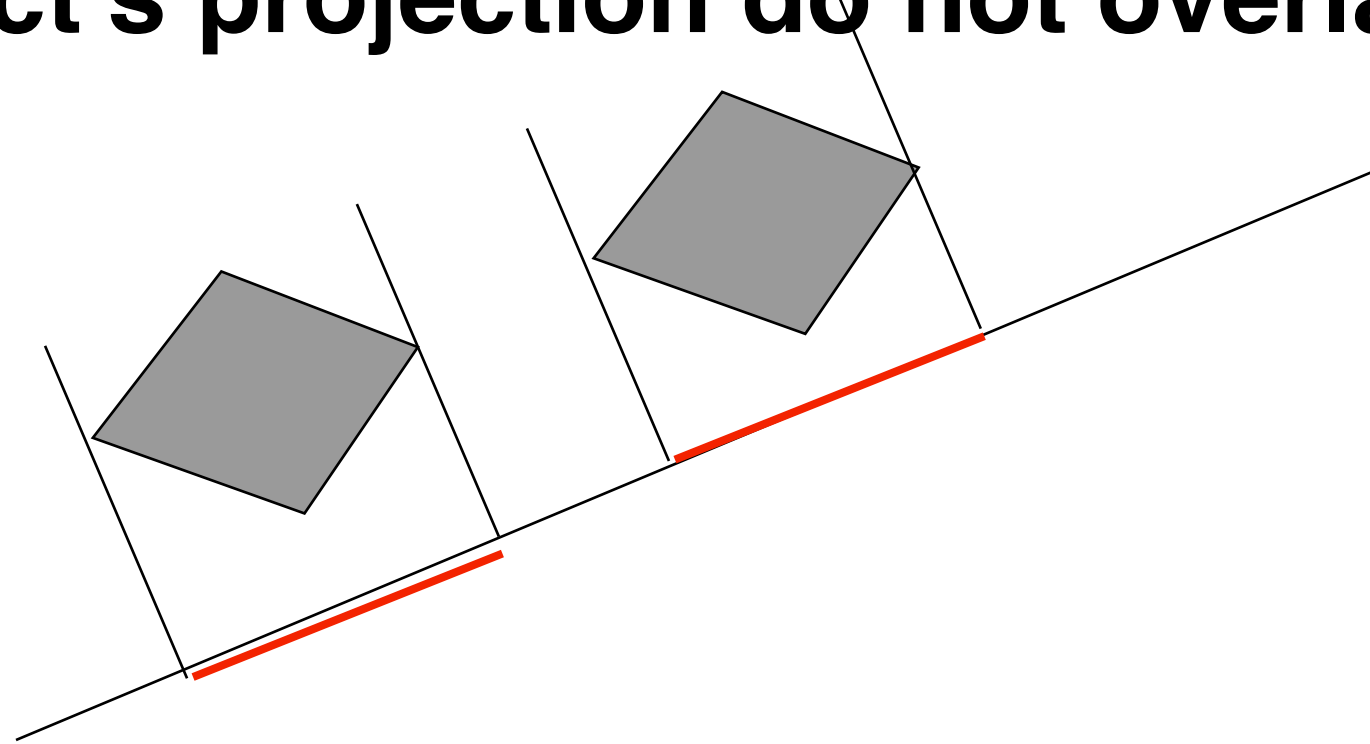
Separating axis theorem (SAT)

**and its application on collision
detection**



Separating axis theorem

Two convex objects do *not* overlap if there exists a line (axis) onto which the two object's projection do not overlap





Really "Separating plane"

Axis = normal vector of separating plane

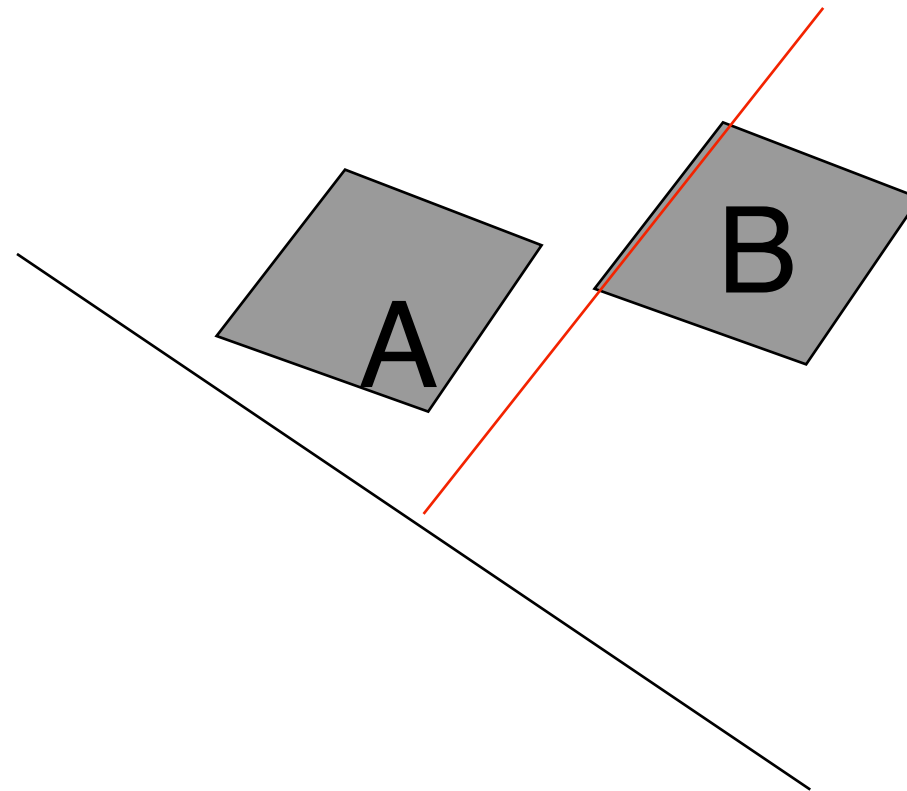
Can you find a plane between them?

**Problem: The number of possible planes/
axes to test are *infinite!***



SAT in practice

For polyhedra models, use the *faces* of the models!

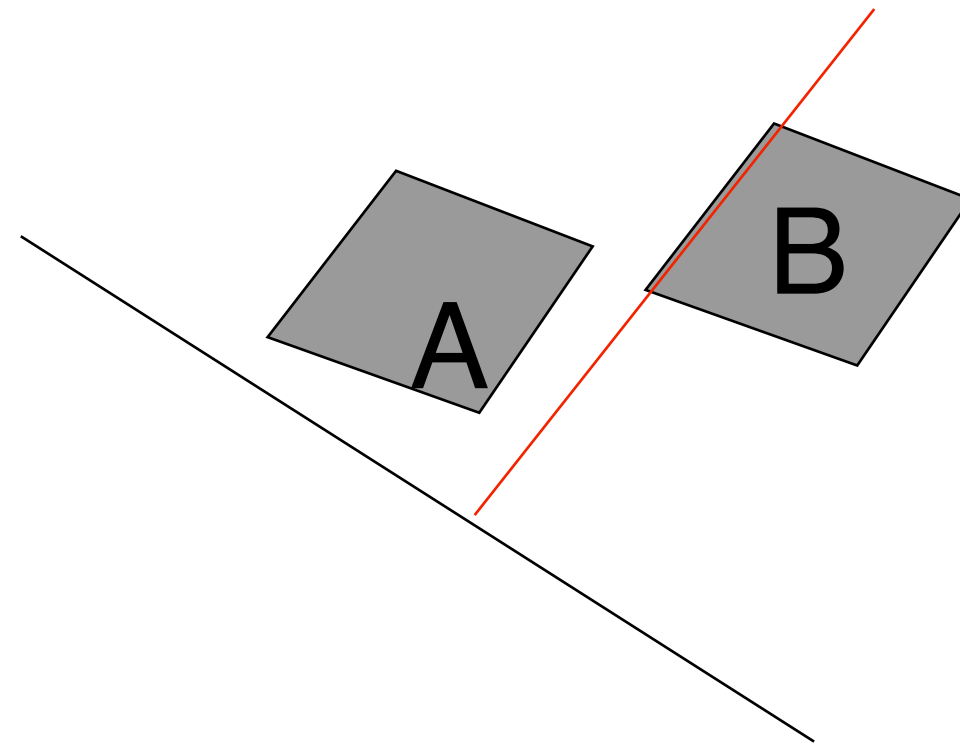




Fairly simple in the 2D case

For every face in B, find plane equation

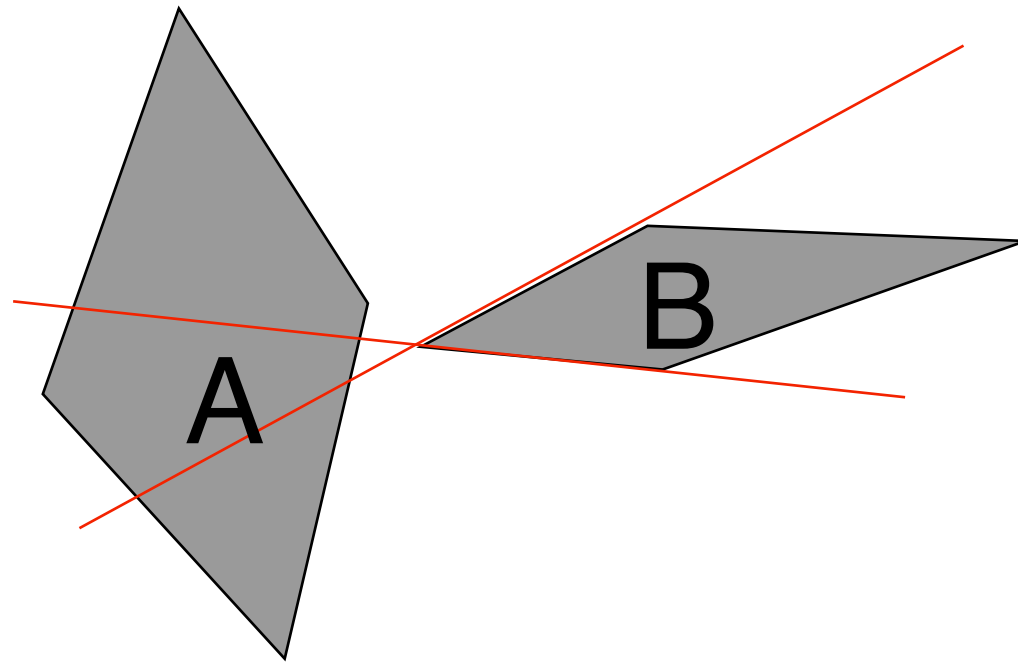
Test all vertices in A with dot product!





But you must go both ways!

There are configurations where one model has no single face that will exclude the other!





2D SAT algorithm

```
for all faces in A
  let a be a vertex in A
  hit = false
  for all vertices b in B
    diff = n · a - n · b
    if diff > 0 then
      hit = true
    if not hit then
      return false

  for all faces in B
    (same algorithm with A and B reversed)

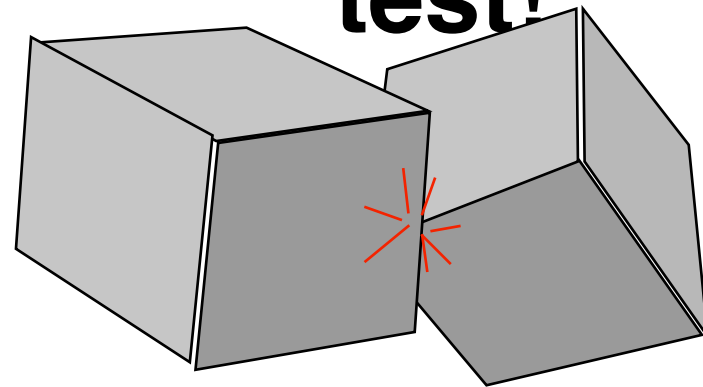
return true
```



SAT in 3D

Not as simple!

Two object meeting at an edge may fail the test!



We must add tests with more planes.



SAT in 3D

We must add tests with more planes. Which ones?

Cross product of an edge in a with an edge in b = additional potential separating plane

All edges in a * all edges in b!

=> Complete SAT on complex 3D objects is *expensive!*

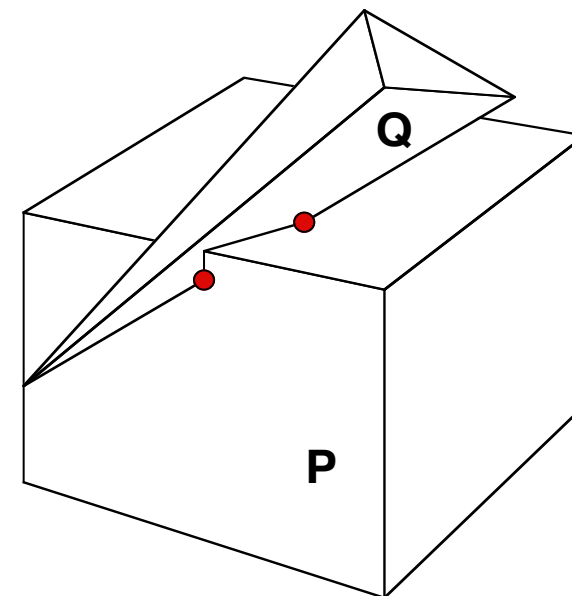
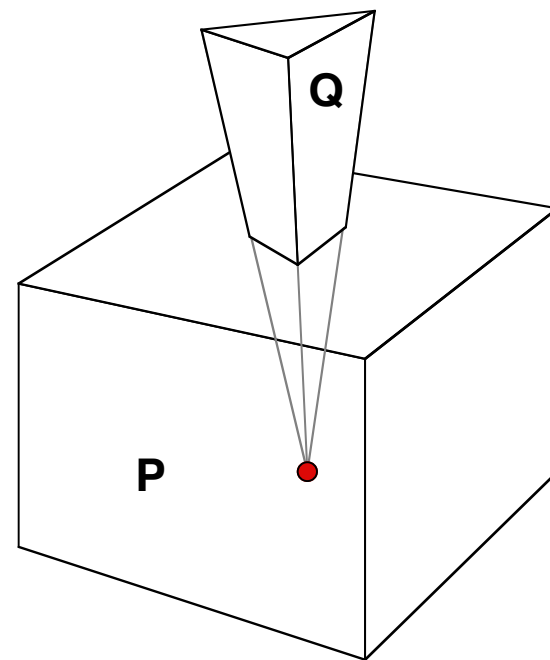


So how can SAT be useful?

- **Simplify. Skip some tests, accept some false hits.**
- **Use SAT on simplified bounding shapes.**
- **Remember separating axis from previous test!**



Containment/intersection





Containment/Intersection

Exact test between two polyhedra, Q and P.

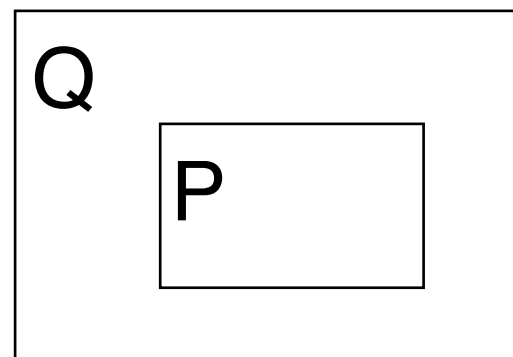
Two tests are needed:

- 1) Is any vertex in Q contained in P or vice versa?**
- 2) Does any edge of Q penetrate a face of P?**



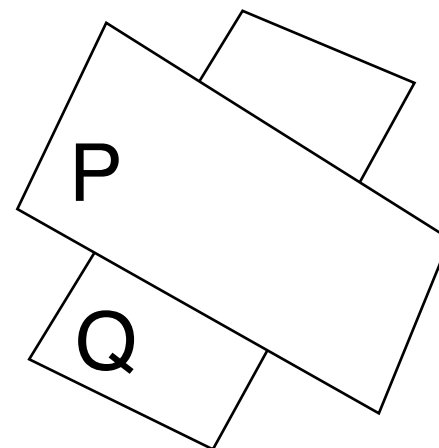
The need for all tests

1) in both directions:



**Q contains points in P,
P contains no points in Q**

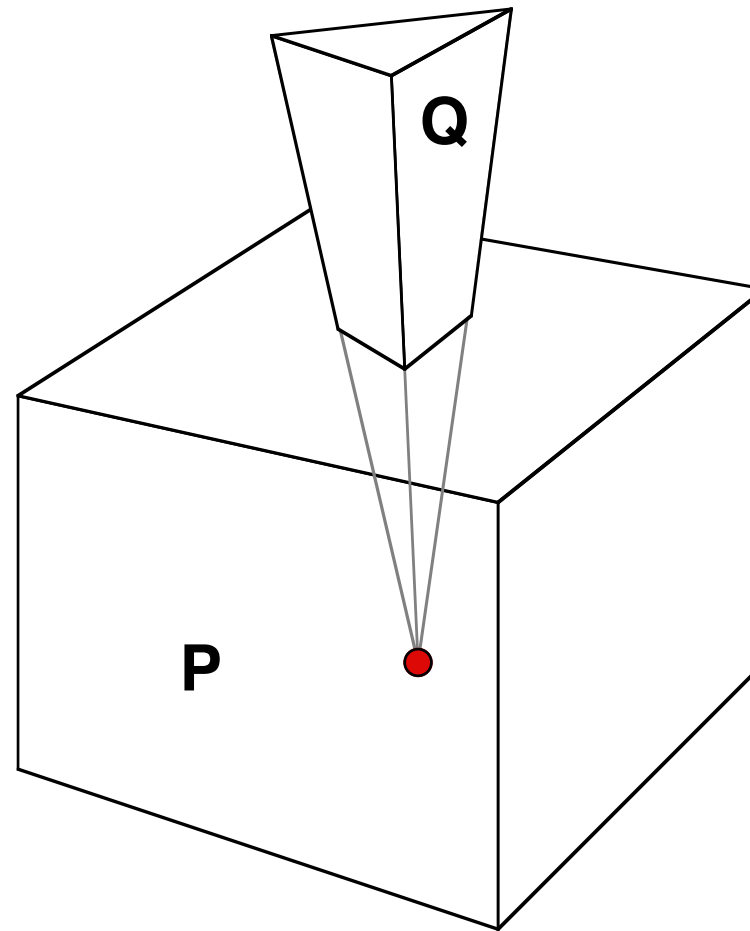
2) when 1) fails:



Neither P nor Q contain points of each other, but still overlap!



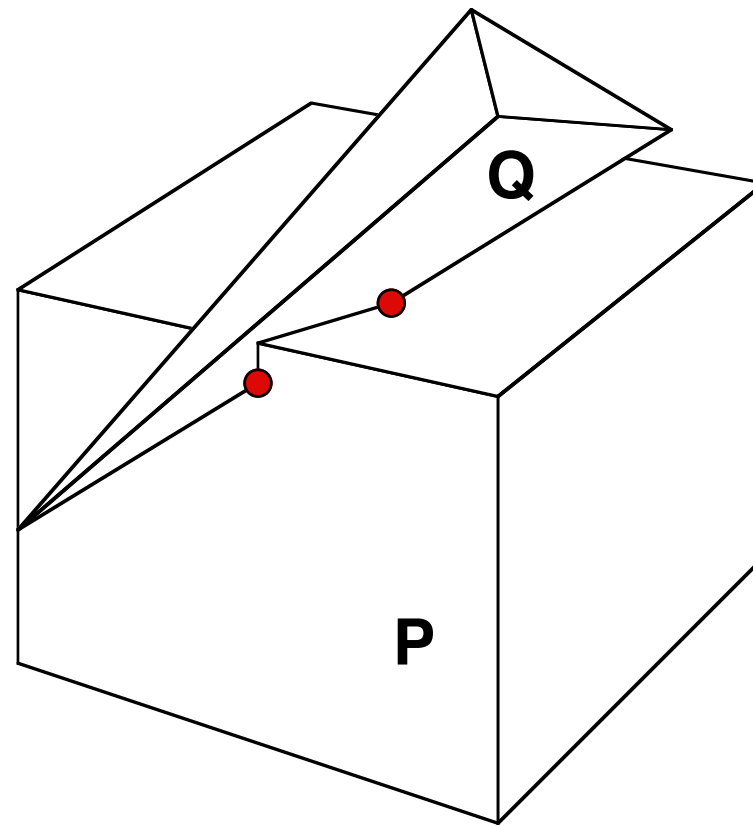
Test 1:



A point in Q is inside P



Test 2:



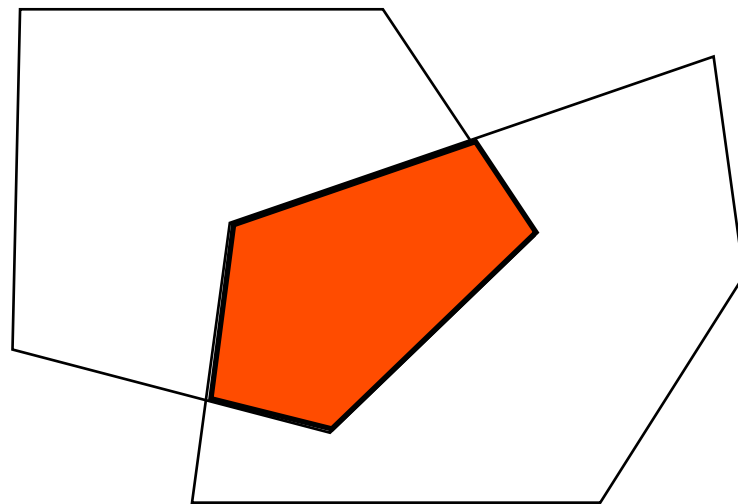
An edge of Q intersects sides of P



Volume intersection

Split P by any plane in Q , repeat. When all planes have been used, the intersection has been calculated.

Can be fairly fast if a good first plane is chosen.





Comparison

Full SAT

$$V_p * P_q + P_p * V_q + E_p * E_q * V_p + E_p * E_q * V_q \quad O(N^3)$$

Containment/intersection:

$$V_p * P_q + P_p * V_q + E_p * P_q + P_p * E_q \quad O(N^2)$$

Volume intersection:

$$\sum k^i P_q \approx P_q \log P_p$$

$$O(N \log N)$$

where $i = 0$ to P_p and $k < 1$.



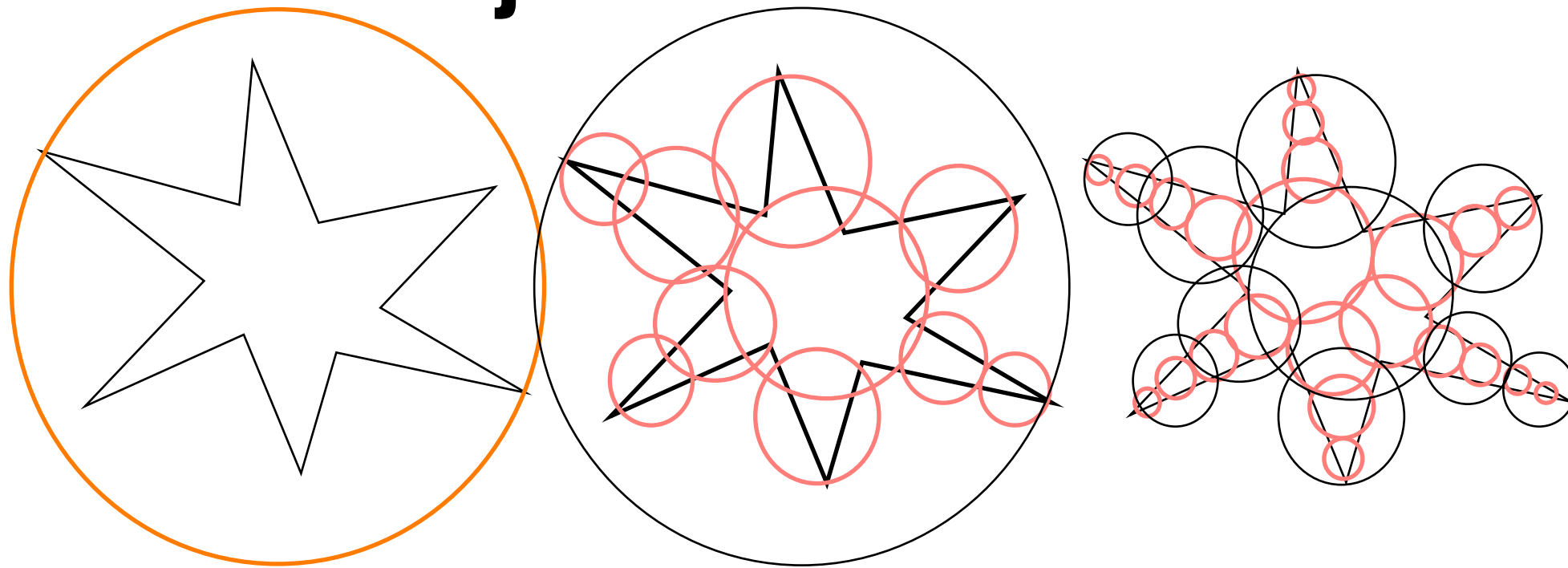
Advanced methods

Volume intersection can sometimes simplify by starting "right".

**Advanced methods take this even longer.
Remember last time, reduce test to very few. Just
test the "right test".**



Single phase algorithms Object hierarchies



**Collision detection is done to the level that
available time permits**